Section 4: Threads and Context Switching

CS162

February 10, 2017

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1 Warmup

1.1 Hello World

What does C print in the following code?

```c
void* identify(void* arg) {
  pid_t pid = getpid();
  printf("My pid is %d\n", pid);
  return NULL;
}

int main() {
  pthread_t thread;
  pthread_create(&thread, NULL, identify, NULL);
  identify(NULL);
  return 0;
}
```

2 Vocabulary

- **thread** - a thread of execution is the smallest unit of sequential instructions that can be scheduled for execution by the operating system. Multiple threads can share the same address space, but each thread independently operates using its own program counter.

- **pthreads** - A POSIX-compliant (standard specified by IEEE) implementation of threads. A `pthread_t` is usually just an alias for “unsigned long int”.

- **pthread_create** - Creates and immediately starts a child thread running in the same address space of the thread that spawned it. The child executes starting from the function specified. Internally, this is implemented by calling the clone syscall.

  ```c
  /* On success, pthread_create() returns 0; on error, it returns an error number, and the contents of *thread are undefined. */
  int pthread_create(pthread_t *thread, const pthread_attr_t *attr, 
      void *(*start_routine) (void *), void *arg);
  ```

- **pthread_join** - Waits for a specific thread to terminate, similar to `waitpid(3)`.

  ```c
  /* On success, pthread_join() returns 0; on error, it returns an error number. */
  int pthread_join(pthread_t thread, void **retval);
  ```

- **pthread_yield** - Equivalent to thread_yield() in Pintos. Causes the calling thread to vacate the CPU and go back into the ready queue without blocking. The calling thread is able to be scheduled again immediately. This is not the same as an interrupt and will succeed in Pintos even if interrupts are disabled.

  ```c
  /* On success, pthread_yield() returns 0; on error, it returns an error number. */
  int pthread_yield(void);
  ```
3 Problems

3.1 Join

What does C print in the following code?
(Hint: There may be zero, one, or multiple answers.)

```c
void *helper(void *arg) {
    printf("HELPER\n");
    return NULL;
}

int main() {
    pthread_t thread;
    pthread_create(&thread, NULL, &helper, NULL);
    pthread_yield();
    printf("MAIN\n");
    return 0;
}
```

How can we modify the code above to always print out "HELPER" followed by "MAIN"?
3.2 Stack Allocation

What does C print in the following code?

```c
void *helper(void *arg) {
    int *num = (int*) arg;
    *num = 2;
    return NULL;
}

int main() {
    int i = 0;
    pthread_t thread;
    pthread_create(&thread, NULL, &helper, &i);
    pthread_join(thread, NULL);
    printf("i is %d\n", i);
    return 0;
}
```

3.3 Heap Allocation

What does C print in the following code?

```c
void *helper(void *arg) {
    char *message = (char *) arg;
    strcpy(message, "I am the child");
    return NULL;
}

int main() {
    char *message = malloc(100);
    strcpy(message, "I am the parent");
    pthread_t thread;
    pthread_create(&thread, NULL, &helper, message);
    pthread_join(thread, NULL);
    printf("%s\n", message);
    return 0;
}
```
3.4 Threads and Processes

What does C print in the following code?
(Hint: There may be zero, one, or multiple answers.)

```c
void *worker(void *arg) {
    int *data = (int *) arg;
    *data = *data + 1;
    printf("Data is %d\n", *data);
    return (void *) 42;
}

int data;
int main() {
    int status;
    data = 0;
    pthread_t thread;

    pid_t pid = fork();
    if (pid == 0) {
        pthread_create(&thread, NULL, &worker, &data);
        pthread_join(thread, NULL);
    } else {
        pthread_create(&thread, NULL, &worker, &data);
        pthread_join(thread, NULL);
        pthread_create(&thread, NULL, &worker, &data);
        pthread_join(thread, NULL);
        wait(&status);
    }
    return 0;
}
```

How would you retrieve the return value of worker? (e.g. "42")
3.5 Context Switching

Refer to the “Pintos Context Switch” section at the end of this discussion worksheet to answer these questions:

How many stacks are involved in a context switch? Identify the purpose of each stack.

The value of `SWITCH_CUR` is 20. The value of `SWITCH_NEXT` is 24. With this information, please draw the stack frame of `switch_threads` for a thread that is about to switch the stack pointer to the next thread’s stack. Your stack frame should include the arguments `cur` and `next`.

In addition to the code inside `switch_threads`, what other actions are required to perform a context switch between 2 user program threads?

In order to perform a context switch, the kernel must copy all of a thread’s registers onto the CPU’s registers. How is the `%eip` (instruction pointer) register copied onto the CPU? Identify which instruction is responsible for this.
3.6 Interrupt Handlers

Refer to the “Pintos Interrupt Handler” section at the end of this discussion worksheet to answer these questions:

What do the instructions `pushal` and `popal` do?

The interrupt service routine (ISR) must run with the kernel’s stack. Why is this the case? And which instruction is responsible for switching the stack pointer to the kernel stack?

The `pushal` instruction pushes 8 values onto the stack (32 bytes). With this information, please draw the stack at the moment when “call intr_handler” is about to be executed.

What is the purpose of the “`pushl %esp`” instruction that is right before ”call intr_handler”?

Inside the `intr_exit` function, what would happen if we reversed the order of the 5 `pop` instructions?
3.7 Pintos Context Switch

```assembly
.globl switch_threads
.func switch_threads
switch_threads:
    # Save caller’s register state.
    # Note that the SVR4 ABI allows us to destroy %eax, %ecx, %edx,
    # but requires us to preserve %ebx, %ebp, %esi, %edi. See
    # [SysV-ABI-386] pages 3-11 and 3-12 for details.
    # This stack frame must match the one set up by thread_create()
    # in size.
    pushl %ebx
    pushl %ebp
    pushl %esi
    pushl %edi

    # Get offsetof (struct thread, stack).
    .globl thread_stack_ofs
    mov thread_stack_ofs, %edx

    # Save current stack pointer to old thread’s stack, if any.
    movl SWITCH_CUR(%esp), %eax
    movl %esp, (%eax,%edx,1)

    # Restore stack pointer from new thread’s stack.
    movl SWITCH_NEXT(%esp), %ecx
    movl (%ecx,%edx,1), %esp

    # Restore caller’s register state.
    popl %edi
    popl %esi
    popl %ebp
    popl %ebx
    ret
.endfunc
```
3.8 Pintos Interrupt Handler

```asm
/*
 * An example of an entry point that would reside in the interrupt vector. This entry point is for interrupt number 0x30.
 */
.func intr30_stub
intr30_stub:
    pushl %ebp /* Frame pointer */
    pushl $0 /* Error code */
    pushl $0x30 /* Interrupt vector number */
    jmp intr_entry
.endfunc

/* Main interrupt entry point. */
.func intr_entry
intr_entry:
    /* Save caller's registers. */
    pushl %ds
    pushl %es
    pushl %fs
    pushl %gs
    pushal

    /* Set up kernel environment. */
    cld /* String instructions go upward. */
    mov $SEL_KDSEG, %eax /* Initialize segment registers. */
    mov %eax, %ds
    mov %eax, %es
    leal 56(%esp), %ebp /* Set up frame pointer. */
    addl $4, %esp

    /* Call interrupt handler. */
    pushl %esp
    .globl intr_handler
    call intr_handler
.endfunc
```
48 /* Interrupt exit.
49
50 Restores the caller’s registers, discards extra data on the
51 stack, and returns to the caller.
52
53 This is a separate function because it is called directly when
54 we launch a new user process (see start_process() in
55 userprog/process.c). */
56 .globl intr_exit
57 .func intr_exit
58 intr_exit:
59 /* Restore caller’s registers. */
60 popal
61 popl %gs
62 popl %fs
63 popl %es
64 popl %ds
65 /* Discard ‘struct intr_frame’ vec_no, error_code,
66 frame_pointer members. */
67 addl $12, %esp
68 /* Return to caller. */
69 iret
70 .endfunc